Language Models

Human Language from a Computational Perspective April 11, 2018

Natural Language Processing

Algorithms that understand or generate human language

Hebrew	English	Japanese	Detect language	•	€.	English	Hebrew	Hungaria	n -	Translate		
mac	hine	transl	ation	תרגום מכונה ⊘								
what is question answering?										Į	<u> </u>	۹
Que infor build lang	stion matior ding sy uage	Answer n retrieve stems th	ing (QA) is a Il and natural l at automatica	comp langua ally ans	uter ige p swe	science process r quest	e discip sing (NL tions p	oline with P), which osed by	hin th ch is ⁄ hun	າe fields of concerned nans in a ກະ	wit atur	th al

computer program in order to create a summary that retains the most important points of the original document. Technologies that can make a coherent summary take into account variables such as length, writing style and syntax. Automatic data summarization is part of machine learning and data mining. The main idea of summarization is to find a representative subset of the data, which contains the information of the entire set. Summarization technologies are used in a large number of sectors in industry today. An example of the use of summarization technology is search engines such as Google. Other examples include document summarization, image collection summarization and video summarization. Document summarization, tries to automatically create a representative summary or abstract of the entire document, by finding the most informative sentences. Similarly, in image summarization the system finds the most representative and important (or salient) images. Similarly, in consumer videos one would want to remove the boring or repetitive scenes, and extract out a much shorter and concise version of the video.

Automatic summarization: reducing text with a computer to retain the most important points.

Statistical Language Model

How likely is each of these sentences? PLEASE MAKE ME A CUP OF COFFEE PLEASE MAKE ME A CUP OF BUTTER PLEASE MAKE ME A CUP OF BOTTLE PLEASE MAKE ME A CUP OF DREAM PLEASE MAKE ME A CUP OF PLEASE

Uses of Language Models

- Typing prediction
- Spelling correction
- Speech recognition
- Many more



Algorithm

Instructions for manipulating data.

Can get parameters as input.

Returns an **output**.



Pseudocode

Notation to describe algorithms.

Not a programming language, but clear enough for humans.

Find the largest number in a list.

$$[3, 1, 4, 16, 0, 2] \rightarrow 16$$

$$[1, 2, 1, 1, 1] \rightarrow 2$$

$$\left[-3, -2, 0, -1\right] \rightarrow 0$$

max(L): m ← L[1] i ← 2 while $i \leq \text{len}(L)$: if L[i] > m: $m \leftarrow L[i]$ i ← i + 1 return m

- L is a list of numbers
- assign the first number to m
- assign 2 to i
- repeat while i is at most len(L)
- b the i'th number is larger than m
- > assign the i'th number to m
- ▷ increase i by 1
- output is the value of m

max(L): Comments $m \leftarrow L[1]$ i ← 2 while $i \leq len(L)$: if L[i] > m: $m \leftarrow L[i]$ $i \leftarrow i + 1$

- repeat while i is at most len(L)
- b the i'th number is larger than m
- assign the i'th number to m
- ▷ increase i by 1
- output is the value of m

Function > L is a list of numbers max $m \leftarrow L[1]$ **definition** \triangleright assign the first number to m i ← 2 ▷ assign 2 to i repeat while i is at most len(L) while $i \leq len(L)$: b the i'th number is larger than m if L[i] > m: assign the i'th number to m $m \leftarrow L[i]$ $i \leftarrow i + 1$ ▷ increase i by 1 output is the value of m return m

Parameters ▷ L is a list of numbers assign the first number to m $m \leftarrow L[1]$ i ← 2 ▷ assign 2 to i while $i \leq \text{len}(L)$: repeat while i is at most len(L) if **L**[i] > m: b the i'th number is larger than m assign the i'th number to m m ← L[i A function can get more than one parameter, but max gets just one

max(L): $\mathbf{m} \leftarrow L[1]$ i – 2 Variables while $i \leq len(L)$: if L[**i**] > **m**: $m \leftarrow L[i]$ i ← i + 1 return m

- repeat while i is at most len(L)
- b the i'th number is larger than m
- assign the i'th number to m
- ▷ increase i by 1
- output is the value of m

L is a list of numbers max(L): assign the first number to m $m \leftarrow L[1]$ i ← 2 Assignment ▷ assign 2 to i repeat while i is at most len(L) while $i \leq len(L)$: b the i'th number is larger than m if L[i] > m: assign the i'th number to m $m \leftarrow L[i]$ i ← i + 1 ▷ increase i by 1 output is the value of m return m

- max(L): i ← 2 while $i \leq len(L)$: if L**[i]** > m: $m \leftarrow L[i]$ $i \leftarrow i + 1$ return m
 - x(L): \triangleright L is a list of numbers $m \leftarrow L[1]$ Indexing $i \leftarrow 2$ \triangleright assign the first number to m \triangleright assign 2 to i
 - repeat while i is at most len(L)
 - b the i'th number is larger than m
 - assign the i'th number to m
 - ▷ increase i by 1
 - output is the value of m

max(L): L is a list of numbers $m \leftarrow L[1]$ Function \triangleright assign the first number to m call i ← 2 ▷ assign 2 to i while $i \leq \text{len}(L)$: repeat while i is at most len(L) b the i'th number is larger than m if L[i] > m: assign the i'th number to m m ← L[i] The function **len** returns the number of elements (length) of a list

max(L): $m \leftarrow L[1]$ $i \leftarrow 2$ Loop while $i \leq \text{len}(L)$: if L[i] > m: $m \leftarrow L[i]$ i ← i + 1

return m

- repeat while i is at most len(L)
- b the i'th number is larger than m
- assign the i'th number to m
- ▷ increase i by 1
- output is the value of m

max(L): $m \leftarrow L[1]$ $i \leftarrow 2$ Condition while $i \leq \text{len}(L)$: if L[i] > m: $m \leftarrow L[i]$ $i \leftarrow i + 1$

return m

- repeat while i is at most len(L)
- b the i'th number is larger than m
- assign the i'th number to m
- ▷ increase i by 1
- output is the value of m

max(L): $m \leftarrow L[1]$ i ← 2 while $i \leq len(L)$: if L[i] > m: $m \leftarrow L[i]$ **Output** $i \leftarrow i + 1$ return m

- repeat while i is at most len(L)
- b the i'th number is larger than m
- assign the i'th number to m
- ▷ increase i by 1
- output is the value of m

- max(L): Indentation → m ← L[1] $i \leftarrow 2$ \checkmark while i \leq len(L): → if L[i] > m: ____>m ← L[i] i ← i + 1
- return m

- L is a list of numbers
 assign the first number to m
 assign 2 to i
- repeat while i is at most len(L)
- b the i'th number is larger than m
- assign the i'th number to m
- ▷ increase i by 1
- output is the value of m

max(L): m ← L[1] i ← 2 while $i \leq \text{len}(L)$: if L[i] > m: $m \leftarrow L[i]$ i ← i + 1 return m

L = [3, 1, 4, 16, 0, 2]m i

max(L): $m \leftarrow L[1]$ i ← 2 while $i \leq len(L)$: if L[i] > m: $m \leftarrow L[i]$ $i \leftarrow i + 1$ return m

L = [3, 1, 4, 16, 0, 2] m i

max(L): m ← L[1] i ← 2 while $i \leq len(L)$: if L[i] > m: $m \leftarrow L[i]$ $i \leftarrow i + 1$ return m

L = [3, 1, 4, 16, 0, 2] m = 3

max(L): $m \leftarrow L[1]$ i ← 2 while $i \leq len(L)$: if L[i] > m: $m \leftarrow L[i]$ $i \leftarrow i + 1$ return m

```
L = [3, 1, 4, 16, 0, 2]
m = 3
i = 2
```

max(L): $m \leftarrow L[1]$ i ← 2 while $i \leq \text{len}(L)$: if L[i] > m: $m \leftarrow L[i]$ $i \leftarrow i + 1$ return m

L = [3, 1, 4, 16, 0, 2]m = 3 i = 2 len(L) = 6 $2 \le 6$

max(L): $m \leftarrow L[1]$ i ← 2 while $i \leq len(L)$: if L[i] > m: $m \leftarrow L[i]$ $i \leftarrow i + 1$ return m

L = [3, 1, 4, 16, 0, 2]m = 3 i = 2 L[i] = L[2] = 1 1 > 3

max(L): $m \leftarrow L[1]$ i ← 2 while $i \leq len(L)$: if L[i] > m: $m \leftarrow L[i]$ i ← i + 1 return m

```
L = [3, 1, 4, 16, 0, 2]
m = 3
i = 3
```

max(L): $m \leftarrow L[1]$ i ← 2 while $i \leq \text{len}(L)$: if L[i] > m: $m \leftarrow L[i]$ $i \leftarrow i + 1$ return m

L = [3, 1, 4, 16, 0, 2]m = 3 i = 3 len(L) = 6 $3 \le 6$

max(L): $m \leftarrow L[1]$ i ← 2 while $i \leq len(L)$: if L[i] > m: $m \leftarrow L[i]$ $i \leftarrow i + 1$ return m

L = [3, 1, 4, 16, 0, 2]m = 3 i = 3 L[i] = L[3] = 4 4 > 3

max(L): $m \leftarrow L[1]$ i ← 2 while $i \leq \text{len}(L)$: if L[i] > m: $m \leftarrow L[i]$ $i \leftarrow i + 1$ return m

L = [3, 1, 4, 16, 0, 2] m = 4 i = 3 L[i] = L[3] = 4

max(L): $m \leftarrow L[1]$ i ← 2 while $i \leq len(L)$: if L[i] > m: $m \leftarrow L[i]$ i ← i + 1 return m

```
L = [3, 1, 4, 16, 0, 2]
m = 4
i = 4
```

max(L): $m \leftarrow L[1]$ i ← 2 while $i \leq \text{len}(L)$: if L[i] > m: $m \leftarrow L[i]$ $i \leftarrow i + 1$ return m

L = [3, 1, 4, 16, 0, 2]m = 4 i = 4 len(L) = 6 4 ≤ 6

max(L): $m \leftarrow L[1]$ i ← 2 while $i \leq len(L)$: if L[i] > m: $m \leftarrow L[i]$ $i \leftarrow i + 1$ return m

L = [3, 1, 4, 16, 0, 2]m = 4 i = 4 L[i] = L[4] = 16 16 > 4

max(L): $m \leftarrow L[1]$ i ← 2 while $i \leq \text{len}(L)$: if L[i] > m: $m \leftarrow L[i]$ $i \leftarrow i + 1$ return m

L = [3, 1, 4, 16, 0, 2] m = 16 i = 4 L[i] = L[4] = 16

max(L): $m \leftarrow L[1]$ i ← 2 while $i \leq len(L)$: if L[i] > m: $m \leftarrow L[i]$ i ← i + 1 return m

L = [3, 1, 4, 16, 0, 2] m = 16 i = 5

max(L): $m \leftarrow L[1]$ i ← 2 while $i \leq \text{len}(L)$: if L[i] > m: $m \leftarrow L[i]$ $i \leftarrow i + 1$ return m

L = [3, 1, 4, 16, 0, 2]m = 16 i = 5 len(L) = 6 $5 \le 6$

max(L): $m \leftarrow L[1]$ i ← 2 while $i \leq len(L)$: if L[i] > m: $m \leftarrow L[i]$ $i \leftarrow i + 1$ return m

L = [3, 1, 4, 16, 0, 2]m = 16 i = 5 L[i] = L[5] = 0 0 > 16
max(L): $m \leftarrow L[1]$ i ← 2 while $i \leq len(L)$: if L[i] > m: $m \leftarrow L[i]$ i ← i + 1 return m

L = [3, 1, 4, 16, 0, 2] m = 16 i = 6

max(L): $m \leftarrow L[1]$ i ← 2 while $i \leq \text{len}(L)$: if L[i] > m: $m \leftarrow L[i]$ $i \leftarrow i + 1$ return m

L = [3, 1, 4, 16, 0, 2]m = 16 i = 6 len(L) = 6 $6 \le 6$

max(L): $m \leftarrow L[1]$ i ← 2 while $i \leq len(L)$: if L[i] > m: $m \leftarrow L[i]$ $i \leftarrow i + 1$ return m

L = [3, 1, 4, 16, 0, 2]m = 16 i = 6 L[i] = L[6] = 2 2 > 16

max(L): $m \leftarrow L[1]$ i ← 2 while $i \leq len(L)$: if L[i] > m: $m \leftarrow L[i]$ i ← i + 1 return m

L = [3, 1, 4, 16, 0, 2] m = 16 i = 7

max(L): $m \leftarrow L[1]$ i ← 2 while $i \leq \text{len}(L)$: if L[i] > m: $m \leftarrow L[i]$ $i \leftarrow i + 1$ return m

L = [3, 1, 4, 16, 0, 2]m = 16 i = 7 len(L) = 67 ≤ 6 🗡

max(L): $m \leftarrow L[1]$ i ← 2 while $i \leq len(L)$: if L[i] > m: $m \leftarrow L[i]$ $i \leftarrow i + 1$

return m

output: 16

Finding index of maximum

Index of largest number in a list.

$$[3, 1, 4, 16, 0, 2] \rightarrow 4$$

 $[1, 2, 1, 1, 1] \rightarrow 2$

$$\left[-3, -2, 0, -1\right] \rightarrow 3$$

Finding index of maximum

argmax(L): a ← 1 i ← 2 while $i \leq \text{len}(L)$: if L[i] > L[a]: a ← i i ← i + 1 return a

- L is a list of numbers
- index of the first element
- index of the second element
- repeat while i is at most len(L)
- i'th number is larger than a'th
- ⊳ assign i to a
- ▷ increase i by 1
- output: index of largest number

Back to language models

Given a list of tokens, predict the most

likely token to follow.

PLEASE MAKE ME A CUP OF <u>TEA</u>

Tokenization

We represent a string:

"I'M LATE!", HE SAID.

As a **list** of tokens:

" I 'M LATE ! " , HE SAID .

Language model algorithm

Predict the next token in the list.

PLEASE MAKE ME A CUP OF \rightarrow COFFEE ONE TWO THREE \rightarrow FOUR WHAT IS YOUR PHONE \rightarrow NUMBER

But the list of tokens is not enough.

We also need to know the language.

Corpora

A text corpus is used to

analyze the distribution of

words.

EARLY CORPUS LINGUISTICS

imputities: It was Busa and Juilland together who worked out much of the foundations of modern corpus linguistics. If their contributions are less than well known, it is angely because corpus linguistics because closely associated with work in English only, and neither Busa nor Juilland worked on corpora of modern English. To review work in English linguistics, we need to comider the last two major groups working on corpus linguistics from the 195th onwards.

1.5.3. Work arising from the study of English grammar

Work on English corpus linguistics started in the early 1960s, when Ouirk (1960) planned and executed the construction of his ambitious Survey of English Usage (sttt). In the same year, Francis and Kucera began work on the now famous Brown corpus, a work which was to take almost two decades to complete." These researchers were in a minority, but they were not universally regarded as peculiar in the field of English language studies, and others followed their lead. For example, in 1975, fourteen years after work began on the Brown corpus, Jan Svartvik started to build on the work of the satu and the Brown corpus to construct the London-Lund corpus. The computer became the mainstay of English corpus linguistics in the 1970s. Svartvik computerised the SEU and, as a consequence, produced what, as Leech (1991: 9) said, was for a long time 'an unmatched resource for the study of spoken English'. With the apprarance in the mid 1990s of large-scale corpora of spontaneous spoken English, such as the spoken section of the British National Corpus, the importance of the London-Lund corpus has faded somewhat. Even so, it is still the only corpus of spontaneous spoken English annotated with prosodic mark-up and, as such, still retains a niche in modern corpus linguistics to this day.

Another side effect of the work of the stut was the training of academics in a tradition of corpus-based approaches to the grammatical analysis of English. Geoffrey Leech was associated with the early days of the sist and went on to Lancater to start a corpus research centre which has given rise to a number of well-known corpus-bailding projects, including perhaps most finously the Lancater-Aolo-Bergen corpus (tos) and more recently the British National Corpus. Sidney Genehaam was also associated with the stut as sometime assitant director of the stri the mid-1980s and founded the International Quirk as director of the stri to the mid-1980s and founded the International Corpus of English project. So the stut was important as it spawned useful corpus resources and trained some linguists who later went on to become pioneers in the field of English corpus linguistics.

The work of Prancis and Kuczer, as well as that of Quirk and his disciples, impired centres of English corpus huiding and exploitation beyond the United Kingdom in addition to Lund. English corpus Inguistics in the tradition of the strust studie grave in Europe throughout the 1970s and 1980s with centres for corpus work being established serons Scandinavia (e.g. Bergen, Gottenburg, Cholo), Weistern Exproy (e.g., Berlin, Chenniatz, Nijmegen) and CORPUS LINGUSTICS FROM THE 1988S TO THE EARLY 1988S

Entern Europe¹⁶ (e.g. Leipzig, Posdam). There is listle doubt that a great deal of the current popularity of corpus linguistics, especially in studies of the English linguage, can be traced to this line of work. However, another related, though somewhat separate, strand of corpus work has been similarly influential in English corpus linguistics over the past forty years. That is the work of the neo-Firthian.

1.5.4. Work by neo-Firthians

3. P., Firth had a colourful life by any definition of the term. He studied in Leeds, picked up an interest in language while on military service in Afghanistan, Africa and India and went on to be professor firstly of English at Lahore and Later at the School of Oriental and African Studies from 1944 (where he started as a senior lecturer in 1938, being appointed professor in 1944). His impact upon English – and more specifically Breitsh – Ingiparities has been notable. He was deeply inflarenced by the work of the anthropologies Brenilaw Mallinowski and the phonetician Daniel Jones: Thip produced a series of writings in the 1930, 1940s and 1950s which were published in a upmope of communication are paramounc?"

The central concept ... is the context of situation. In that context are the human participant or participants, what they say, what is going on. The phonetician can find his phonetic context, and the grammarian and the lexicographer theirs.

Firth's agenda dominated much of British linguistics for the best part of a generation. As stated, its penetration beyond the United Kingdom. In America in particular it cannot be stild that Firth's views ever constituted a dominant paradigm of linguistic research. Indeed, some of the most trenchart criticism of the Firthian approach to language can be found in the writing of American Inguists, e.g. Langendeen (1968), hough it is possible to find even stronger critical voices raised by others, e.g. Lyson (1968), Firth's place in corpus linguistics is assumed, however, largely because he stated (Firth, 1977- 29), that 'Attested language..., uhy recorded is in the focus of attention for the linguist' and used some terminology which is used to this day in corpus languistics. Journal we solve stores that we will call neo-Firthian linguists, such as Halliday, Hoey and Sinchair, to work in the tradition the statishibade.

On the terminology side, his term collacation is in use in modern corpus linguistics to this day (see section 3.4.4 for example). However, the popularity of that term can most easily be understood in the context of later corpus linguists, such as Sinchär, using the term. Collocation as a concept has a history

Counts table

To represent token	,	775
	THE	630
counts, we map strings	•	392
to numbers.	"	345
	AND	339
Some ready-made counts:	А	337
books.google.com/ngrams	ТО	277

Algorithm to count words

Count all words in a tokenized corpus

and return a table of counts.

[I, AM, SAM, .,	
SAM, I, AM., -	\rightarrow
I, DO, NOT, LIKE,	
green, eggs, and, ham, .]	

Ι	3	LIKE	1
AM	2	GREEN	1
SAM	2	EGGS	1
	3	AND	1
DO	1	НАМ	1
NOT	1		

Algorithm to count words

- count(L): C1 ← [0] i ← 1 while $i \leq \text{len}(L)$: t ← L[i] $C1[t] \leftarrow C1[t] + 1$ i ← i + 1 return C1
- L is a list of tokens
 create a table of zeros
- assign 1 to i
- repeat while i is at most len(L)
- get token at position i
- ▷ increase count for t by 1
- increase i by 1
- output is the counts table

Algorithm to count words

count(L): Using a word as an index to a table while $i \leq \text{len}(L)$: t ← L[i] $C1[t] \leftarrow C1[t] + 1$ $i \leftarrow i + 1$ return C1

L is a list of tokens

- create a table of zeros
- ▷ assign 1 to i
- repeat while i is at most len(L)
- get token at position i
- ▷ increase count for t by 1
- increase i by 1
- output is the counts table

Word counts

Example counts from Alice's Adventures in Wonderland (1866) by Lewis Carroll



Unigram Language Model

Easiest: always predict	3	775
the most frequent token:	THE	630
I wish I <mark>,</mark>		392
C1 =	"	345
(Unigram counts)	AND	339
C1[,] = 775	A	337
С1[тне] = 630	ТО	277

Bigram counts

We can also count	, THE	530
bigrams (pairs of words)	AND THE	320
		89
C2 =	SHE SAID	65
C2[and, the] = 320		
C2[she, said] = 65	Ι 'м	20
	I do	10

Bigram Language Model

Look only at the last	I	'м	20
token to predict the next:	I	DO	10
I wish <mark>I 'm</mark>	I	'LL	10
C2[I, ·] =	I	'VE	10
(Bigram counts starting	I	SHOULD	8
with I)	I	MUST	7
	Ι	THINK	7

C2[I, 'м] = 20 C2[I, do] = 10

Trigram Language Model

Look at the **two** last

tokens to predict the next:

I wish I could

C3[wish, I, ·] =

(Trigram counts starting

with wish I)

C3[wish, I, could] = 20 C3[wish, I, над] = 10

wish I could20wish I had10

n-gram Language Model

Look at the *n* – 1 last

tokens to predict the next:

I wish I could

C4[I, wish, I, ⋅] =

(4-gram counts starting

with I wish I):

C3[I, wish, I, could] = 2 C3[I, wish, I, наd] = 2

I wish I could 2 I wish I had 2

Algorithm to count *n*-grams

count(L, n): $[0] \rightarrow O$ i ← 1 while $i \leq \text{len}(L) - n + 1$: $T \leftarrow L[i, ..., i + n - 1]$ $C[T] \leftarrow C[T] + 1$ i ← i + 1 return C

- L: list of tokens, n: a number
- create a table of zeros
- ⊳ assign 1 to i
- repeat while i is at most len(L) n + 1
- get n tokens starting at i
- increase count for T by 1
- ▷ increase i by 1
- output is the counts table

Algorithm to count *n*-grams

- count(L, n): **Getting several** elements from a list while $i \leq len(L) - n + 1$: T ← L[i, ..., i + n − 1] $C[T] \leftarrow C[T] + 1$ $i \leftarrow i + 1$ return C
- L: list of tokens, n: a number
- create a table of zeros
- ⊳ assign 1 to i
- ▷ repeat while i is at most len(L) n + 1
- get n tokens starting at i
- ▷ increase count for T by 1
- ▷ increase i by 1
- output is the counts table

Algorithm to count *n*-grams

count(L, n):

Using an *n*-gram as an index to a table

while $i \leq len(L) - n + 1$:

$$T \leftarrow L[i, ..., i + n - 7]$$
$$C[T] \leftarrow C[T] + 1$$

return C

- L: list of tokens, n: a number
- create a table of zeros
- ⊳ assign 1 to i
- repeat while i is at most len(L) n + 1
- get n tokens starting at i
- ▷ increase count for T by 1
- ▷ increase i by 1
- output is the counts table

Unigram algorithm

unigram(L, C1): return argmax(C1)

▷ L: tokens, C1: unigram counts

b token with highest count

Unigram algorithm

unigram(L, C1): returnargmax(C1)

Function call

- L: tokens, C1: unigram counts
- token with highest count

Unigram algorithm

unigram(L, C1): return argmax(C1)

L: tokens, C1: unigram counts

token with highest count

Ignores L and always predicts the same word...

Bigram algorithm

 $\begin{aligned} \textbf{bigram}(L, C2): \\ k \leftarrow len(L) \\ t \leftarrow L[k] \\ return argmax(C2[t, \cdot]) \end{aligned}$

- ▷ L: tokens, C2: bigram counts
- ▷ length of L
- last token in L
 - bigram with highest count,
- ▷ among the bigrams starting with t

Bigram algorithm

```
bigram(L, C2):

k \leftarrow len(L)

t \leftarrow L[k]

return argmax(C2[t, ·])
```

Getting part of the table

- L: tokens, C2: bigram counts
- ▷ length of L
- ▷ last token in L
 - bigram with highest count,
- among the bigrams starting with t

Trigram algorithm

trigram(L, C3): $k \leftarrow len(L)$ $T \leftarrow L[k - 1, k]$ return argmax(C3[T, ·])

- ▷ L: tokens, C3: trigram counts
- ▷ length of L
- Iast two tokens in L
- trigram with highest count,
- ▷ among the trigrams starting with T

General *n*-gram algorithm

 $\begin{array}{l} \textbf{ngram}(L, n, Cn):\\ k \leftarrow len(L)\\ T \leftarrow L[k - n + 2, ..., k]\\ \textbf{return argmax}(Cn[T, \cdot]) \end{array}$

▷ L: tokens, Cn: *n*-gram counts

▷ length of L

last n – 1 tokens in L

n-gram with highest count,

among the n-grams starting with T

This can replace **unigram**, **bigram** and **trigram** algorithms: just use n=1, n=2 or n=3

Text prediction algorithm

predict(L, n, Cn, m):

 $P \leftarrow L$

- L: tokens, Cn: n-gram counts,
- m: total wanted number of words
- start with words given as input

n-gram models comparison

Unigram	, , , , , , , , , , , , , , , , , , , ,
Bigram	Then she went on it had been running about in her head $!$ the garden , who was now , for some of them $!$
Trigram	ALL OF A GOOD DEAL FRIGHTENED AT THE TOP OF HER SISTER , WHO WAS GENTLY BRUSHING AWAY SOME DEAD LEAVES THAT HAD FALLEN INTO A TREE A FEW MINUTES , IT WAS THE WHITE RABBIT , WHO WAS NOW ABOUT TWO FEET HIGH .
4-gram	THE FIRST THING I 'VE GOT TO DO , SO ALICE SOON BEGAN TALKING TO HERSELF . `` DINAH 'LL MISS ME VERY MUCH TO-NIGHT , I SHOULD THINK ! " (DINAH WAS THE CAT
5-gram	AND SO IT WAS INDEED ! SHE WAS NOW ONLY TEN INCHES HIGH , AND HER FACE BRIGHTENED UP AT THE THOUGHT THAT SHE WAS NOW ABOUT TWO FEET HIGH AND WAS GOING ON SHRINKING RAPIDLY .

Back-off

n-gram models quickly become too **sparse**.

Whenever I wish I

does not occur in Alice in Wonderland: cannot use 5-gram.

If no match is found, use a smaller *n*:

To predict the next token, **back-off** to 4-grams:

Whenever **I wish I could**



Trigram with Backoff to Bigram

trigram-backoff-bigram(L, C2, C3): > L: tokens, $k \leftarrow len(L) \triangleright C2$: bigram counts, C3: trigram counts if C3[L[k - 1, k], \cdot]) is empty : ▷ not found return argmax(C2[L[k], ·]) use bigram else: trigram found return argmax(C3[L[k − 1, k], ·]) buse trigram
Trigram with Backoff to Bigram

trigram-backoff-bigram(L, C2, C3): > L: tokens, $k \leftarrow len(L) \triangleright C2$: bigram counts, C3: trigram counts if C3[L[k - 1, k], \cdot]) is empty : ▷ not found return argmax(C2[L[k], ·]) use bigram else: trigram found return argmax(C3[L[k - 1, k], ·]) ▷ use trigram if/else condition

Trigram with Full Backoff

trigram-backoff(L, C): \triangleright L is a list of tokens, $k \leftarrow len(L)$ \triangleright C is the list [C1, C2, C3]: i ← 3 unigram, bigram, trigram counts while $C[i][L[k - i + 2, ..., k], \cdot]$ is empty: $i \leftarrow i - 1$ ▷ *i*-gram not found, try *i* – 1 return argmax(C[i][L[k - i + 2, ..., k], ·])

References

- Google Ngram Viewer: <u>books.google.com/ngrams</u>
- Alice's Adventures in Wonderland on Wikisource:

en.wikisource.org/wiki/Alice's_Adventures_in_Wonderland_(1866)

• *n*-grams: <u>en.wikipedia.org/wiki/N-gram</u>



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